

Universität Freiburg Institut für Informatik Michael Meier Fang Wei Georges-Köhler Allee, Geb. 51 D-79110 Freiburg Tel. (0761) 203-8126 Tel. (0761) 203-8125

## Foundations of Query Languages Summer semester 2010 May 4, 2010

# 3. Exercise Set: FOL and Conjunctive Queries

### Exercise 1

Consider the following database schema Prof(PName), Course(CName), GivesCourse(PName,CName), Student(SName), Attends(SName,CName), Grades(PName,SName). Write down in the relational calculus:

- a) What professors offer no courses?
- b) What students are graded by a professor who offer no courses?
- c) What pairs of students attend a common course?
- d) What professors offer a course and also grades all students in that course?

#### Exercise 2

Consider the following pairs of conjunctive queries and decide for each pair  $q_i$ ,  $q'_i$  if  $q_i \sqsubseteq q'_i$ ,  $q'_i \sqsubseteq q_i$ , and  $q_i \equiv q'_i$  holds. If such relationships hold provide the corresponding containment mappings. Otherwise, give a database that provides a counterexample

a)  $q_1: \mathtt{ans}(X,Y) \leftarrow \mathtt{R}(X,Z), \mathtt{R}(Z,T), \mathtt{S}(T,Y) \text{ and } q'_1: \mathtt{ans}(X,Z) \leftarrow \mathtt{R}(X,X), \mathtt{S}(X,Z)$ 

- b)  $q_2$ : ans(X)  $\leftarrow R(X,Y), S(Y,Z), S(Y',Z')$  and  $q'_2$ : ans(Y)  $\leftarrow S(A,B), R(Y,A), R(Y',A)$
- c)  $q_3$ : ans(U,Z)  $\leftarrow R(U,V), R(X,Y), S(Y,Z), S(V,X) \text{ and } q'_3$ : ans(U,V)  $\leftarrow R(Y,U), R(U,X), S(U,V), S(X,Y)$

#### Exercise 3

Consider the following pairs of conjunctive queries and decide if  $q_i \sqsubseteq q'_i$ ,  $q'_i \sqsubseteq q_i$ , and  $q_i \equiv q'_i$  hold using the method of the canonical instance.

a)  $q_1$ : ans(X)  $\leftarrow R(X,Y,X), R(X,Z,Y), S(Y,X)$  and  $q'_1$ : ans(X)  $\leftarrow R(X,Y,Z), S(Y,Z)$ b)  $q_2$ : ans(X)  $\leftarrow R(X,Y), R(Y,Z), R(Z,X)$  and  $q'_2$ : ans(X)  $\leftarrow R(X,Y), R(Y,Z), R(Z,U), R(U,V)$ 

#### Exercise 4

Let Q, Q' be conjunctive queries. Prove or disprove the following statements:

- a) Every conjunctive query Q is satisfiable, i.e. for every conjunctive query Q there is a database instance I such that  $Q(I) \neq \emptyset$ .
- b)  $Q(I) \subseteq Q'(I)$  for all finite database instances I iff it holds that  $Q(J) \subseteq Q'(J)$  for all (possibly infinite) database instances J.

Due by: May 12, 2010 before the tutorial starts.